# Rationals vs Byzantine Players in Committee-Based Blockchains

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#### **BLOCKCHAIN**

Potentially unbounded set of processes that communicate in a network through message passing

Distributed ledger

**Tamper-resistant** 

Build in an append only manner





### CONSENSUS

### Termination

Every non-faulty process eventually decides some value

### Agreement

If there is a non-faulty process that decides a value B, then eventually all the Non-faulty processes decide B

## Validity

A decided value is valid, it satisfies the predefined predicate



**COMMITTEE-BASED BLOCKCHAINS** 

Committee

# **Dissecting Tendermint**



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# HotStuff: BFT Consensus with Linearity and Responsiveness

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block to be appended

• Committees are rewarded for their work

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## **RATIONAL BEHAVIOUR & SYSTEM MODEL**

**Q.** Are the consensus properties (Termination and Validity) guaranteed with the presence of rational participants ?

Ordered set of n processes/players

Messages are signed and signatures cannot be forged Processes cannot lie about who created a message

#### Synchronous communication

Messages cannot be lost

Following the BAR Model[1], participants are either

- o Rational
- o Byzantine

Assumption: There are more rational processes than Byzantine

[1] A. S. Aiyer, L. Alvisi, A. Clement, M. Dahlin, J.-P. Martin, and C. Porth, 'BAR fault tolerance for cooperative services', in *Proceedings of the 20th ACM Symposium on Operating Systems Principles (SOSP'05)*, 2005, pp. 45–58.

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**OBJECTIVES** 



$$(R * \mathbb{1}_{(\sigma_i^{\text{send}}(H_i^T)=1)} * \mathbb{1}_{(\text{block accepted at }T)} - \kappa \mathbb{1}_{(\text{invalid block accepted})}) - \sum_{t=1}^T \left( c_{check} \mathbb{1}_{\sigma_i^{\text{check}}(h_i^t)=1)} + c_{send} \mathbb{1}_{(\sigma_i^{\text{send}}(H_i^t)=1)} \right)$$

 $\kappa > R > c_{check} > c_{send}$ 



#### **EXAMPLES OF EXECUTION** 3 MESSAGES REQUIRED





#### Perfect Bayesian equilibrium

- 1. Choose actions maximizing their objective function,
- 2. Rationally anticipate the strategies of the others, and
- **3.** Draw rational inferences from what they observe, using their expectations about the strategies of the others and Bayes law, whenever it applies.





#### **CONSENSUS AGAINST RATIONALS**

#### Q. Are the consensus properties (Termination and Validity) guaranteed with the presence of rational participants ?

*n* is the total number of processes *v* is the minimum number of required messages for block's production *f* is the number of Byzantine processes (f>1)



➤ When f ≥ v, in equilibrium, all rational participants send a message without checking validity

Termination holds

Validity is not guaranteed



When f < v, there exists an equilibrium where all rational participants do not check block's validity, nor send a message





- When f < v, if the cost of producing an invalid block is "high enough", there is an equilibrium where there is always a valid block produced
  - For a process i, if i ≤ n-v+f+1 then i checks block's validity and sends a message iff the block is valid
  - For a process i, if i > n-v+f+n (i ≤ n) i sends a message without checking block's validity

*Termination* holds *Validity* holds



**CONCLUSIONS & PERSPECTIVES** 

Analyse of rational behaviour in Committee-based Blockchains against Byzantine processes

**Extend the current work with more settings** 

# Merci ! Thank You !







#### EXAMPLE OF EXECUTION (1/2) 3 MESSAGES REQUIRED

