## Large Poisson Games

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Based on Matías Núñez, "Approval Voting in Large Electorates", in Handbook on Approval Voting, Springer-Verlag, 2010.

## Position of the problem

Example: Voting

- Most of times, my action has no impact at all!
- If the population is fixed and other players act deterministically, then generally, any possible action is a best response for me (because it does not matter!).
- This is a difficulty to define equilibria.

If there is uncertainty on the population of players:

- There is always a small probability that a pivotal situation arises, where my action matters.
- I can choose my action based on these very unlikely events.

Poisson games is just a possible model to introduce population uncertainty (but it offers important practical advantages for mathematical tractability).

## Overview

Principle:

1. Each player believes that the other players will use some strategy.
2. She compute the (unlikely) events where her action makes a difference.
3. She choose her strategy as a best response to this analysis.

Equilibria are defined as fixed points of this process.

## References

- Roger Myerson (2000). Large Poisson games. Journal of Economic Theory, 94, 7-45.
- Proves existence of equilibria
- Magnitude Theorem
- Roger Myerson (2002). Comparison of scoring rules in Poisson voting games. Journal of Economic Theory, 103, 219-251.
- Dual Magnitude Theorem
- Matías Núñez (2010). Condorcet consistency of approval voting: A counter example on large Poisson games. Journal of Theoretical Politics, 22(1), 64-84.
- Magnitude Equivalence Theorem
- Matías Núñez (2010). Approval Voting in Large Electorates, in Handbook on Approval Voting, Springer-Verlag, 2010.
- Reader's digest of all the above (+ alternative models)


## Plan

Model

Direct probability calculation
Magnitude Theorem (Myerson, 2000)

Dual Magnitude Theorem (Myerson, 2002)

Magnitude Equivalence Theorem (Núñez 2010)

Conclusion

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Conclusion

## Our running example: an election in Approval voting

Approval voting:

- Each voter votes for (= approves of) any number of candidates.
- The score of a candidate is the number of votes she receives.
- The candidate with highest score is declared the winner.
- In case of tie, the winner is chosen uniformly at random among the candidates with highest score.


## Drawing the population of voters ( $=$ players )

- Expected number of voters: $n$.
- Actual number of voters: $N \sim \mathcal{P}(n)$ (Poisson distribution with mean $n$ ).

$$
\mathbb{P}[N=k]=e^{-n} \frac{n^{k}}{k!}
$$

## Drawing the types of the voters

Each voter's type is independently drawn.

| Type $t$ | $t_{1}$ | $t_{2}$ | $t_{3}$ |
| :--- | :---: | :---: | :---: |
|  | $\alpha$ | $\beta$ | $\gamma$ |
| Preference ranking | $\beta$ | $\alpha$ | $\alpha$ |
|  | $\gamma$ | $\gamma$ | $\beta$ |
| Type distribution $r(t)$ | 0.1 | 0.6 | 0.3 |
| Actual number of voters $N(t)$ | $\sim \mathcal{P}(0.1 n)$ | $\sim \mathcal{P}(0.6 n)$ | $\sim \mathcal{P}(0.3 n)$ |

- Random variables $N(t)$ are independent.
- Types also have utilities (not written in the above table). For example, $u_{t_{1}}(\alpha)>u_{t_{1}}(\beta)>u_{t_{1}}(\gamma)$.


## Choosing ballots (= actions)

For example, consider this strategy function: $\left\{\begin{array}{l}\sigma\left(\alpha \mid t_{1}\right)=1 \\ \sigma\left(\alpha \beta \mid t_{2}\right)=1 \\ \sigma\left(\gamma \mid t_{3}\right)=1\end{array}\right.$.

- $\sigma$ is generally not given. The issue will precisely be to find a $\sigma$ that yields an equilibrium.
Then we have:

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| Ballot distribution $\tau(c)$ | 0.1 | 0.6 | 0.3 |
| Actual number of ballots $X(c)$ | $\sim \mathcal{P}(0.1 n)$ | $\sim \mathcal{P}(0.6 n)$ | $\sim \mathcal{P}(0.3 n)$ |

- Random variables $X(c)$ are independent.
- What happens if $\sigma\left(\alpha \beta \mid t_{1}\right)=1$ instead?


## Computing scores

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

Scores:

| Candidate $\kappa$ | $\alpha$ | $\beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| Score distribution $\rho(\kappa)$ | 0.7 | 0.6 | 0.3 |
| Actual score $S(\kappa)$ | $\sim \mathcal{P}(0.7 n)$ | $\sim \mathcal{P}(0.6 n)$ | $\sim \mathcal{P}(0.3 n)$ |

- Are random variables $S(\kappa)$ independent?
- The winner is candidate $\alpha$.


## Advantages of the Poisson model

Common public information $=$ Environmental equivalence
From the point of view of any voter, the number of other voters follows $\mathcal{P}(n)$, the number of other voters with type $t$ follows $\mathcal{P}(\operatorname{nr}(t))$, etc. Hence all voters live in the same environment, which is the same as seen by an external observer.

Independence of actions
The number $X(c)$ of voters who choose a given ballot is independent from the number of voters who choose another ballot.

## Large Poisson games

We consider a sequence of Poisson games, parametrized by the expected number of players $n$.
Limit properties when $n \rightarrow \infty$ ?

## Plan

## Model

Direct probability calculation
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Conclusion

## Direct probability calculation: Ex. 1

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

Event $B_{n}$ : there is no ballot $\gamma$.

$$
\mathbb{P}\left[X \in B_{n}\right]=\mathbb{P}[X(\gamma)=0]=e^{-n \tau(\gamma)} \frac{(n \tau(\gamma))^{0}}{0!}=e^{-n \tau(\gamma)}
$$

The magnitude of $B=\left(B_{n}\right)_{n \in \mathbb{N}}$ is defined as the coefficient in the exponent:

$$
\mu(B)=\lim _{n \rightarrow \infty} \frac{1}{n} \log \mathbb{P}\left[X \in B_{n}\right]=-\tau(\gamma)
$$

- A magnitude is always $\leq 0$.
- When $n \rightarrow \infty$, it is unlikely that there is no ballot $\gamma: e^{-0.3 n}$.
- But it is infinitely less likely that there is no ballot $\alpha \beta: e^{-0.6 n}$.
- Similarly, it is infinitely more likely that there is no ballot $\alpha: e^{-0.1 n}$.


## Direct probability calculation: Ex. 2

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

Event $B_{n}$ : there are just as many ballots $\alpha \beta$ as $\gamma$.

$$
\begin{aligned}
& \mathbb{P}\left[X \in B_{n}\right]=\sum_{k=0}^{\infty} \mathbb{P}[X(\alpha \beta)=k \text { and } X(\gamma)=k] \\
&=\sum_{k=0}^{\infty} \mathbb{P}[X(\alpha \beta)=k] \cdot \mathbb{P}[X(\gamma)=k] \\
&=e^{-n(\tau(\alpha \beta)+\tau(\gamma))} \sum_{k=0}^{\infty} \frac{\left(n^{2} \tau(\alpha \beta) \tau(\gamma)\right)^{k}}{(k!)^{2}} \\
&=e^{-n(\tau(\alpha \beta)+\tau(\gamma))} l_{0}(2 n \sqrt{\tau(\alpha \beta) \tau(\gamma)}) \\
&=\exp (-n(\tau(\alpha \beta)+\tau(\gamma))+2 n \sqrt{\tau(\alpha \beta) \tau(\gamma)}+o(n)) \\
& \mu(B)=-\tau(\alpha \beta)-\tau(\gamma)+2 \sqrt{\tau(\alpha \beta) \tau(\gamma)}=-(\sqrt{\tau(\alpha \beta)}-\sqrt{\tau(\gamma)})^{2}
\end{aligned}
$$

## Plan

## Model

Direct probability calculation

Magnitude Theorem (Myerson, 2000)
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Conclusion

## Magnitude of a sequence of points

$b_{n}$ sequence of points. For each kind of ballot $c$, it specifies how many ballots $b_{n}(c)$.

$$
\begin{aligned}
\mu(b) & =\lim _{n \rightarrow \infty} \frac{1}{n} \log \mathbb{P}\left[X=b_{n}\right] \\
& =\lim _{n \rightarrow \infty} \frac{1}{n} \log \prod_{c \in \mathcal{C}} \mathbb{P}\left[X(c)=b_{n}(c)\right] \\
& =\lim _{n \rightarrow \infty} \frac{1}{n} \sum_{c \in \mathcal{C}} \log \left(e^{-n \tau_{n}(c)} \frac{\left(n \tau_{n}(c)\right)^{b_{n}(c)}}{b_{n}(c)!}\right) \\
& =\lim _{n \rightarrow \infty} \sum_{c \in \mathcal{C}} \tau_{n}(c)\left(\frac{b_{n}(c)}{n \tau_{n}(c)}\left(1-\log \frac{b_{n}(c)}{n \tau_{n}(c)}\right)-1\right)
\end{aligned}
$$

## Offset

Hence we have:

$$
\mu(b)=\lim _{n \rightarrow \infty} \sum_{c \in \mathcal{C}} \tau_{n}(c) \psi\left(\frac{b_{n}(c)}{n \tau_{n}(c)}\right),
$$

where $\psi(x)=x(1-\log x)-1$ and $\psi(0)=-1$.


We define the offset of ballot $c$ in this sequence of points $b_{n}$ as:

$$
\phi_{c}=\lim _{n \rightarrow \infty} \frac{b_{n}(c)}{n \tau_{n}(c)} .
$$

It is the limit ratio between the number of actual ballots $c$ in this particular sequence of points and what would be expected in general.

## Magnitude Theorem

Let $\left(B_{n}\right)$ be a sequence of outcomes whose magnitude is defined. Then:

$$
\begin{aligned}
\mu(B) & =\lim _{n \rightarrow \infty} \frac{1}{n} \log \mathbb{P}\left[X \in B_{n}\right] \\
& =\lim _{n \rightarrow \infty} \max _{b_{n} \in B_{n}} \frac{1}{n} \log \mathbb{P}\left[X=b_{n}\right]
\end{aligned}
$$

I.e. if $b_{n}$ is defined as the most probable point in $B_{n}$, then $\mu(B)=\mu(b)$. In practice:

$$
\mu(B)=\lim _{n \rightarrow \infty} \max _{b_{n} \in B_{n}} \sum_{c \in \mathcal{C}} \tau_{n}(c) \psi\left(\frac{b_{n}(c)}{n \tau_{n}(c)}\right)
$$

N.B.: We define the offset of ballot $c$ in $B$ as its offset in $b$. [In fact, this offset is the same in any sequence of points extracted from $\left(B_{n}\right)$ that has the same magnitude.]

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

$B_{n}$ : there is no ballot $\gamma$.
Any point $b_{n} \in B_{n}$ is of the form: $X(\alpha)=k, X(\alpha \beta)=k^{\prime}, X(\gamma)=0$.

$$
\begin{aligned}
\mu(B) & =\lim _{n \rightarrow \infty} \max _{b_{n} \in B_{n}} \sum_{c \in \mathcal{C}} \tau_{n}(c) \psi\left(\frac{b_{n}(c)}{n \tau_{n}(c)}\right) \\
& =\lim _{n \rightarrow \infty} \max _{k, k^{\prime} \in \mathbb{N}} \tau(\alpha) \underbrace{\psi\left(\frac{k}{n \tau(\alpha)}\right)}_{\lim \max =0}+\tau(\alpha \beta) \underbrace{\psi\left(\frac{k^{\prime}}{n \tau(\alpha \beta)}\right)}_{\lim \max =0}+\tau(\gamma) \underbrace{\psi\left(\frac{0}{n \tau(\gamma)}\right)}_{=-1} \\
& =-\tau(\gamma)
\end{aligned}
$$

Moreover: $\phi_{\alpha}=\phi_{\alpha \beta}=1$ and $\phi_{\gamma}=0$.

Magnitude Theorem: Ex. 2

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

$B_{n}$ : there are just as many ballots $\alpha \beta$ as $\gamma$.

$$
\begin{aligned}
\mu(B) & =\lim _{n \rightarrow \infty} \max _{b_{n} \in B_{n}} \sum_{c \in \mathcal{C}} \tau_{n}(c) \psi\left(\frac{b_{n}(c)}{n \tau_{n}(c)}\right) \\
& =\lim _{n \rightarrow \infty} \max _{k, k^{\prime} \in \mathbb{N}} \tau(\alpha) \underbrace{\psi\left(\frac{k}{\tau(\alpha) n}\right)}_{\lim \max =0}+\tau(\alpha \beta) \psi\left(\frac{k^{\prime}}{\tau(\alpha \beta) n}\right)+\tau(\gamma) \psi\left(\frac{k^{\prime}}{\tau(\gamma) n}\right) \\
& =\lim _{n \rightarrow \infty} \max _{k^{\prime} \in \mathbb{N}} \tau(\alpha \beta) \psi\left(\frac{k^{\prime}}{\tau(\alpha \beta) n}\right)+\tau(\gamma) \psi\left(\frac{k^{\prime}}{\tau(\gamma) n}\right) \\
& =\max _{x \geq 0} \tau(\alpha \beta) \psi\left(\frac{x}{\tau(\alpha \beta)}\right)+\tau(\gamma) \psi\left(\frac{x}{\tau(\gamma)}\right) \\
& =-(\sqrt{\tau(\alpha \beta)}-\sqrt{\tau(\gamma)})^{2}
\end{aligned}
$$

Magnitude Theorem: Ex. 2 (continued)

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

And the offsets:

$$
\begin{cases}\phi_{\alpha}=1 & \Rightarrow b_{n}(\alpha) \sim 0.1 n \\ \phi_{\alpha \beta}=\sqrt{\tau(\gamma) / \tau(\alpha \beta)} & \Rightarrow b_{n}(\alpha \beta) \sim n \sqrt{\tau(\alpha \beta) \tau(\gamma)} \simeq 0.42 n \\ \phi_{\gamma}=\sqrt{\tau(\alpha \beta) / \tau(\gamma)} & \Rightarrow b_{n}(\gamma) \sim n \sqrt{\tau(\alpha \beta) \tau(\gamma)} \simeq 0.42 n\end{cases}
$$

Remark: the total number of voters in $b_{n}$ is not $n$.

## Myerson 2000: bonus tracks (explained with the hands)

## Offset theorem

Assume you have studied $\left(B_{n}\right)_{n \in \mathbb{N}}$, e.g., $B_{n}: X(\alpha \beta)=X(\gamma)$.
Consider a "finite translation", e.g. $B_{n}^{\prime}: X(\alpha \beta)=X(\gamma)+1$.
The offset theorem gives a easy way to compute:

$$
\lim _{n \rightarrow \infty} \frac{\mathbb{P}\left(B_{n}^{\prime}\right)}{\mathbb{P}\left(B_{n}\right)}
$$

[In particular, $B$ and $B^{\prime}$ have the same magnitude.]

## Hyperplane theorem

Quite technical, but the main consequence is:
For events whose probability does not tend to $\mathbf{0}$, you can approximate by a normal distribution: $X(c) \sim \mathcal{N}(n \tau(c), \sqrt{n \tau(c)})$.

## Plan

## Model

Direct probability calculation

Magnitude Theorem (Myerson, 2000)

Dual Magnitude Theorem (Myerson, 2002)

## Magnitude Equivalence Theorem (Núñez 2010)

## Conclusion

## DMT

Let $B$ be an outcome defined by a finite set of linear inequalities:

$$
\left\{\begin{array}{c}
a_{1}(c) X(c)+a_{1}\left(c^{\prime}\right) X\left(c^{\prime}\right)+\ldots \geq 0 \\
\vdots \\
a_{K}(c) X(c)+a_{K}\left(c^{\prime}\right) X\left(c^{\prime}\right)+\ldots \geq 0
\end{array}\right.
$$

Suppose that $\lambda_{1}, \ldots, \lambda_{K} \geq 0$ is an argmin of:

$$
F(\lambda)=\sum_{c \in \mathcal{C}} \tau(c)\left[\exp \left(\sum_{k \leq K} \lambda_{k} a_{k}(c)\right)-1\right]
$$

Then $\mu(B)=F(\lambda)$. Moreover, for any ballot $c$ :

$$
\phi_{c}=\exp \left(\sum_{k \leq K} \lambda_{k} a_{k}(c)\right)
$$

$B_{n}$ : there is no ballot $\gamma$.
Only one constraint:

$$
\begin{gathered}
0 X(\alpha)+0 X(\alpha \beta)-1 X(\gamma) \geq 0 . \\
F(\lambda)=\tau(\alpha) \underbrace{\left[e^{0 \lambda_{1}}-1\right]}_{=0}+\tau(\alpha \beta) \underbrace{\left[e^{0 \lambda_{1}}-1\right]}_{=0}+\tau(\gamma) \underbrace{\left[e^{-1 \lambda_{1}}-1\right]}_{\rightarrow-1 \text { if } \lambda_{1} \rightarrow \infty} \\
\mu(B)=-\tau(\gamma) \\
\left\{\begin{array}{l}
\phi_{\alpha}=e^{0 \lambda_{1}}=1 \\
\phi_{\alpha \beta}=e^{0 \lambda_{1}}=1 \\
\phi_{\gamma}=e^{-1 \lambda_{1}}=0
\end{array}\right.
\end{gathered}
$$

## DMT: Ex. 2

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

$B_{n}$ : there are just as many ballots $\alpha \beta$ as $\gamma$.
Constraints:

$$
\left\{\begin{array}{l}
0 X(\alpha)+1 X(\alpha \beta)-1 X(\gamma) \geq 0 \\
0 X(\alpha)-1 X(\alpha \beta)+1 X(\gamma) \geq 0
\end{array}\right.
$$

$$
F(\lambda)=\tau(\alpha) \underbrace{\left[e^{0 \lambda_{1}+0 \lambda_{2}}-1\right]}_{=0}+\tau(\alpha \beta)\left[e^{1 \lambda_{1}-1 \lambda_{2}}-1\right]+\tau(\gamma)\left[e^{-1 \lambda_{1}+1 \lambda_{2}}-1\right]
$$

For $e^{\lambda_{1}-\lambda_{2}}=\sqrt{\tau(\gamma) / \tau(\alpha \beta)}$, we obtain:

$$
\mu(B)=2 \sqrt{\tau(\alpha \beta) \tau(\gamma)}-\tau(\alpha \beta)-\tau(\gamma)=-(\sqrt{\tau(\alpha \beta)}-\sqrt{\tau(\gamma)})^{2}
$$

And we could deduce the offsets immediately.

## DMT: About the proof

- In Myerson 2002, the proof is very short.
- But it is a bit obscure where the function $F(\lambda)$ comes from.
- I will try to provide a clue about this!

Kuhn-Tucker conditions: one constraint

$$
\begin{array}{ll}
\max & f(x) \\
\text { s.t. } & a_{1}(\alpha) x_{\alpha}+a_{1}(\beta) x_{\beta} \geq 0
\end{array}
$$

## Example:

$$
-x_{\alpha}+x_{\beta} \geq 0
$$

If $\vec{x}^{*}$ is in the interior of the cone:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=\overrightarrow{0}
$$



## Kuhn-Tucker conditions: one constraint

## $\max f(x)$

s.t. $\quad a_{1}(\alpha) x_{\alpha}+a_{1}(\beta) x_{\beta} \geq 0$

## Example:

$-x_{\alpha}+x_{\beta} \geq 0$

If $\vec{x}^{*}$ is in the interior of the cone:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=\overrightarrow{0}
$$

If $\vec{x}^{*}$ is on the frontier of the cone:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{1} \overrightarrow{a_{1}} \quad\left(\text { with } \lambda_{1} \geq 0\right)
$$



## Kuhn-Tucker conditions: one constraint

## $\max f(x)$

s.t. $\quad a_{1}(\alpha) x_{\alpha}+a_{1}(\beta) x_{\beta} \geq 0$

## Example:

$$
-x_{\alpha}+x_{\beta} \geq 0
$$

If $\vec{x}^{*}$ is in the interior of the cone:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=\overrightarrow{0}
$$

If $\vec{x}^{*}$ is on the frontier of the cone:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{1} \overrightarrow{a_{1}} \quad\left(\text { with } \lambda_{1} \geq 0\right)
$$

Anyway, the second condition is met.
Moreover, if $\lambda_{1}>0$, then $\overrightarrow{a_{1}} \cdot \vec{x}=0$.


## Kuhn-Tucker conditions: several constraints

## $\max f(x)$

s.t. $\quad a_{1}(\alpha) x_{\alpha}+a_{1}(\beta) x_{\beta} \geq 0$

$$
a_{2}(\alpha) x_{\alpha}+a_{2}(\beta) x_{\beta} \geq 0
$$

$\vec{x}^{*}$ interior: $\vec{\nabla} f\left(\vec{x}^{*}\right)=\overrightarrow{0}$.
$\vec{x}^{*}$ on first frontier: $\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{1} \overrightarrow{a_{1}}$.
$\vec{x}^{*}$ on second frontier: $\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{2} \overrightarrow{a_{2}}$.

## Example:

$$
\begin{aligned}
-x_{\alpha}+x_{\beta} & \geq 0 \\
2 x_{\alpha}-0.25 x_{\beta} & \geq 0
\end{aligned}
$$



## Kuhn-Tucker conditions: several constraints

$\max f(x)$
s.t.

$$
\begin{aligned}
& a_{1}(\alpha) x_{\alpha}+a_{1}(\beta) x_{\beta} \geq 0 \\
& a_{2}(\alpha) x_{\alpha}+a_{2}(\beta) x_{\beta} \geq 0
\end{aligned}
$$

$\vec{x}^{*}$ interior: $\vec{\nabla} f\left(\vec{x}^{*}\right)=\overrightarrow{0}$.
$\vec{x}^{*}$ on first frontier: $\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{1} \overrightarrow{a_{1}}$.
$\vec{x}^{*}$ on second frontier: $\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{2} \overrightarrow{a_{2}}$.
If $\vec{x}^{*}$ is on the intersection of frontiers:

$$
\vec{\nabla} f\left(\vec{x}^{*}\right)=-\lambda_{1} \overrightarrow{a_{1}}-\lambda_{2} \overrightarrow{a_{2}} \quad\left(\text { with } \lambda_{k} \geq 0\right)
$$

In all cases, the above condition is met.
Moreover, if $\lambda_{k}>0$, then $\overrightarrow{a_{k}} \cdot \vec{x}=0$.

## Example:

$$
\begin{aligned}
-x_{\alpha}+x_{\beta} & \geq 0 \\
2 x_{\alpha}-0.25 x_{\beta} & \geq 0
\end{aligned}
$$



## Applying Kuhn-Tucker conditions in our case

$$
\begin{array}{ll}
\max & f(\phi)=\tau(\alpha) \psi\left(\phi_{\alpha}\right)+\tau(\beta) \psi\left(\phi_{\beta}\right) \\
\mathrm{s.t.} & a_{1}(\alpha) \tau(\alpha) \phi_{\alpha}+a_{1}(\beta) \tau(\beta) \phi_{\beta} \geq 0 \\
& a_{2}(\alpha) \tau(\alpha) \phi_{\alpha}+a_{2}(\beta) \tau(\beta) \phi_{\beta} \geq 0
\end{array}
$$

Condition on $\vec{\nabla} f$ :

$$
\begin{aligned}
\binom{-\tau(\alpha) \log \phi_{\alpha}}{-\tau(\beta) \log \phi_{\beta}} & =-\lambda_{1}\binom{a_{1}(\alpha) \tau(\alpha)}{a_{1}(\beta) \tau(\beta)}-\lambda_{2}\binom{a_{2}(\alpha) \tau(\alpha)}{a_{2}(\beta) \tau(\beta)} \\
\phi_{\alpha} & \left.=e^{\lambda_{1} a_{1}(\alpha)+\lambda_{2} a_{2}(\alpha)} \quad \text { (resp. } \beta\right)
\end{aligned}
$$

When $\lambda_{k}>0$, then the point is on the corresponding frontier. For example:

$$
\begin{gathered}
a_{1}(\alpha) \tau(\alpha) e^{\lambda_{1} a_{1}(\alpha)+\lambda_{2} a_{2}(\alpha)}+a_{1}(\beta) \tau(\beta) e^{\lambda_{1} a_{1}(\beta)+\lambda_{2} a_{2}(\beta)}=0 \\
\frac{\partial}{\partial \lambda_{1}}\left[\tau(\alpha) e^{\lambda_{1} a_{1}(\alpha)+\lambda_{2} a_{2}(\alpha)}+\tau(\beta) e^{\lambda_{1} a_{1}(\beta)+\lambda_{2} a_{2}(\beta)}\right]=0
\end{gathered}
$$

## Applying Kuhn-Tucker conditions in our case

$$
\begin{array}{ll}
\max & f(\phi)=\tau(\alpha) \psi\left(\phi_{\alpha}\right)+\tau(\beta) \psi\left(\phi_{\beta}\right) \\
\text { s.t. } & a_{1}(\alpha) \tau(\alpha) \phi_{\alpha}+a_{1}(\beta) \tau(\beta) \phi_{\beta} \geq 0 \\
& a_{2}(\alpha) \tau(\alpha) \phi_{\alpha}+a_{2}(\beta) \tau(\beta) \phi_{\beta} \geq 0
\end{array}
$$

Condition on $\vec{\nabla} f$ :

$$
\begin{aligned}
\binom{-\tau(\alpha) \log \phi_{\alpha}}{-\tau(\beta) \log \phi_{\beta}} & =-\lambda_{1}\binom{a_{1}(\alpha) \tau(\alpha)}{a_{1}(\beta) \tau(\beta)}-\lambda_{2}\binom{a_{2}(\alpha) \tau(\alpha)}{a_{2}(\beta) \tau(\beta)} \\
\phi_{\alpha} & \left.=e^{\lambda_{1} a_{1}(\alpha)+\lambda_{2} a_{2}(\alpha)} \quad \text { (resp. } \beta\right)
\end{aligned}
$$

When $\lambda_{k}>0$, then the point is on the corresponding frontier. For example:

$$
\begin{gathered}
a_{1}(\alpha) \tau(\alpha) e^{\lambda_{1} a_{1}(\alpha)+\lambda_{2} a_{2}(\alpha)}+a_{1}(\beta) \tau(\beta) e^{\lambda_{1} a_{1}(\beta)+\lambda_{2} a_{2}(\beta)}=0 \\
\frac{\partial}{\partial \lambda_{1}} F(\lambda)=0
\end{gathered}
$$

## Terminological remarks

- Dual cone: cone generated by the $\vec{a}_{k}$.
- $\lambda_{k}$ : Lagrange multiplier associated to constraint $k$.
- $F(\lambda)$ is the Lagrange function (or Lagrangian) of the original optimization problem. It plays the role of objective function in the dual problem.


## Plan

## Model

Direct probability calculation

## Magnitude Theorem (Myerson, 2000)

Dual Magnitude Theorem (Myerson, 2002)

Magnitude Equivalence Theorem (Núñez 2010)

## Conclusion

Rigorous definition of a pivot situation for subset $Y$ of candidates:

$$
\left\{\begin{array}{l}
\forall \kappa \in Y, S(\kappa) \geq \max S-1 \\
\forall \kappa \notin Y, S(\kappa) \leq \max S-2
\end{array}\right.
$$

These constraints are not linear!
We have $\mu[\operatorname{pivot}(Y)]=\mu(B)$, where outcome $B$ is defined by:

$$
\begin{cases}\forall \kappa, \kappa^{\prime} \in Y: & S(\kappa)=S\left(\kappa^{\prime}\right) \\ \forall \kappa \in Y, \kappa^{\prime} \notin Y: & S(\kappa) \geq S\left(\kappa^{\prime}\right)\end{cases}
$$

These constraints are linear! Hence we can apply DMT.

MET example: pivot $\alpha$ vs $\beta$

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

MET:

$$
\left\{\begin{array}{lrl}
S(\alpha) \geq S(\beta): & 1 X(\alpha) & \geq 0 \\
S(\beta) \geq S(\alpha): & -1 X(\alpha) & \geq 0 \\
S(\alpha) \geq S(\gamma): & 1 X(\alpha)+1 X(\alpha \beta)-1 X(\gamma) \geq 0
\end{array}\right.
$$

DMT:

$$
F(\lambda)=\tau(\alpha)\left[e^{1 \lambda_{1}-1 \lambda_{2}+1 \lambda_{3}}-1\right]+\tau(\alpha \beta)\left[e^{1 \lambda_{3}}-1\right]+\tau(\gamma)\left[e^{-1 \lambda_{3}}-1\right]
$$

Minimization: $\lambda_{1}-\lambda_{2} \rightarrow-\infty$ and $\lambda_{3}=0$ (we cannot do better because $\lambda_{3} \geq 0$ ).
We obtain:

$$
\mu[\operatorname{pivot}(\alpha, \beta)]=-\tau(\alpha)
$$

## End of the example

| Ballot $c$ | $\alpha$ | $\alpha \beta$ | $\gamma$ |
| :--- | :---: | :---: | :---: |
| $\tau(c)$ | 0.1 | 0.6 | 0.3 |

$$
\left\{\begin{array}{l}
\mu[\operatorname{pivot}(\alpha, \beta)]=-\tau(\alpha)=-0.1 \\
\mu[\operatorname{pivot}(\alpha, \gamma)]=-(\sqrt{\tau(\alpha)+\tau(\alpha \beta)}-\sqrt{\tau(\gamma)})^{2} \simeq-0.08 \\
\mu[\operatorname{pivot}(\beta, \gamma)]=-\tau(\alpha)-(\sqrt{\tau(\alpha \beta)}-\sqrt{\tau(\gamma)})^{2} \simeq-0.15 \\
\quad \Rightarrow \mu[\operatorname{pivot}(\alpha, \gamma)]>\mu[\operatorname{pivot}(\alpha, \beta)]>\mu[\operatorname{pivot}(\beta, \gamma)]
\end{array}\right.
$$

| Type $t$ | $t_{1}$ | $t_{2}$ | $t_{3}$ |
| :--- | :---: | :---: | :---: |
| Preference ranking | $\alpha$ | $\beta$ | $\gamma$ |
|  | $\beta$ | $\alpha$ | $\alpha$ |
|  | $\gamma$ | $\gamma$ | $\beta$ |
| Best response | $\alpha$ | $\alpha \beta$ | $\gamma$ |

$\Rightarrow$ It is an equilibrium!

## Plan

```
Model
Direct probability calculation
Magnitude Theorem (Myerson, 2000)
Dual Magnitude Theorem (Myerson, 2002)
Magnitude Equivalence Theorem (Núñez 2010)
```

Conclusion

## Large Poisson Games

- Models with variable population generate uncertainty, even if players have deterministic strategies.
- Players may base their strategy on very unlikely events.
- Poisson games $\Rightarrow$ environmental equivalence $\Rightarrow$ easier to handle!
- Large Poisson games $\Rightarrow$ reason in terms of magnitudes (instead of probabilities).
- Magnitude Theorem, DMT and MET gives practical tools to compute magnitudes.


## Thank you!

